

# Calculating Szeged Index and Revised Szeged Index by Using Adjacency Matrix

Hajar Alimorad

Department of Mathematics, Jahrom University, P.O. Box: 74135-111, Iran

Received: 9 July 2022

Accepted: 17 August 2022

DOI: 10.30473/ijac.2022.64726.1240

## Abstract

Topological indices are graph invariants used in theoretical chemistry to encode molecules for the design of chemical compounds with given physicochemical properties or given pharmacological and biological activities. Szeged index ( $Sz(G)$ ) and Revised Szeged index ( $Sz^*(G)$ ) of molecular graph identify some distance properties for graphs. In computational chemistry and graph theory, ( $Sz(G)$ ) and ( $Sz^*(G)$ ) were more available to test the characteristics of chemical molecular structures, and thus widely used in chemical applications. In this paper, a simple algorithm is presented for constructing the distance matrix. This algorithm is designed for calculation of  $Sz(G)$  and  $Sz^*(G)$ .

## Keywords

Adjacency Matrix; Connected Graph; Distance Matrix; Revised Szeged Index; Szeged Index.

## 1. INTRODUCTION

Suppose  $G = (V, E)$  be a simple and connected graph, with vertex set  $V$  and edge set  $E$ . Let  $n = |V|$  and  $m = |E|$  denote the order and the size of  $G$ .

The distance between two vertices  $u$  and  $v$  in graph is the number of edges in the shortest path and is denoted by  $d(u, v)$ .

Let  $e = uv \in E$  and define the partition, with respect to  $e$ ,  $\{N_u(e), N_v(e), N_0(e)\}$  of the vertices of  $G$  as follows [1]:

$$N_u(e) = \{w \in V : d(u, w) < d(v, w)\},$$

$$N_v(e) = \{w \in V : d(v, w) < d(u, w)\},$$

$$N_0(e) = \{w \in V : d(v, w) = d(u, w)\}.$$

Let  $n_u(e)$ ,  $n_v(e)$  and  $n_0(e)$  denote the number of vertices in  $N_u(e)$ ,  $N_v(e)$  and  $N_0(e)$ , respectively. Szeged index [2, 3] of  $G$  is defined by

$$Sz = Sz(G) = \sum_{e=uv \in E} n_u(e) \cdot n_v(e).$$

The definition of the Szeged index does not take into account the vertices at equal distance to  $u$  and  $v$ . Szeged star index (or revised Szeged index) of  $G$  is defined by [4, 5]:

$$Sz^* = Sz^*(G) = \sum_{e=uv \in E} \left( n_u(e) + \frac{n_0(e)}{2} \right) \cdot \left( n_v(e) + \frac{n_0(e)}{2} \right).$$

Some properties and applications of Szeged index and revised Szeged index have been reported in [1, 6, 7, 8, 9, 10].

This paper is organized in the following way:

After introduction, the new algorithm for obtaining  $Sz(G)$  and  $Sz^*(G)$  is represented in Section 2. In Section 3, effectiveness of the proposed approach is verified by solving numerical examples. Finally, conclusions are discussed in the last section.

## 2. STATEMENT OF THE PROBLEM

The standard distance matrix or the vertex-distance matrix (or the minimum path matrix) of a vertex-labeled connected graph  $G$ , denoted by  $D$ , is a real symmetric  $n \times n$  matrix whose elements are defined as:

$$D(i, j) = \begin{cases} l(i, j) & \text{if } i \neq j, \\ 0 & \text{otherwise,} \end{cases}$$

where  $l(i, j)$  is the length of the shortest path, i. e., the minimum number of edges, between vertices  $i$  and  $j$  in  $G$ .

This matrix has been used to generate a number of topological indices, e.g., Balaban index [11, 12], Wiener index [13, 14], multiplicative Wiener index [15, 16, 17] and distance-sum index [18]. An efficient algorithm is available for computing the vertex-distance matrix of any graph in [19].

A common way to obtain Matrix  $D$  is to use different powers of adjacency Matrix  $A$ .

The elements of Matrix  $A$  give the connections between vertices. Powers of adjacency matrix are concatenating walks. The  $ij$ -th entry of the  $k$ -th power of  $A$  counts the number of walks of length  $k$  from vertex  $i$  to vertex  $j$ , not paths (a walk can repeat vertices, while a path cannot). So, to

create a distance matrix, we need to iteratively power adjacency Matrix  $A$ , and as soon as a  $ij$  –  $th$  element is non-zero, we have to assign the distance  $k$  in distance matrix [20].

```

A = input('Enetr the adjacency matrix');
[n, n] = size(A);
D = NaN(n);
B = A;
k = 1;
while any(isnan(D(:)))
    (Check for new walks and assign distance)
    D(B > 0 & isnan(D)) = k;
    (Iteration)
    k = k + 1;
    B = B * A;
end;
(D(i, i) = 2,    in this algorithm)
for i = 1 : n
    D(i, i) = 0;
end
    
```

In all the mentioned methods, matrix multiplication is used [19, 20].

In this paper, a simple and practical method for calculating Matrix  $D$ , without the need for matrix multiplication will be presented. The advantage of this method is that it can be used for finite graphs with large dimensions and no matrix storage space is necessary.

Matrix  $D$  is symmetric with zero diagonal elements. Therefore, in the proposed algorithm, only the elements above diagonal are calculated. Then, the matrix can be obtained by considering symmetry.

Algorithm:

1. Enter adjacency Matrix  $A$ ,
2. Consider

$$D(i, j) = A(i, j)$$

In this step, all paths with lenght of one are defined.

3. Now starting from  $h = 1$  and until all the values above diagonal of Matrix  $D$  become non-zero

```

for k = 1:n
    for i = k + 1:n
        for j = i + 1:n
            if (A(i, j) = 1 & D(k, i) = h & D(k, j) = 0)
                D(k, j) = h + 1;
                h = h + 1;
            end
        end
    end
end
    
```

Thus, for each  $h$ , all paths of length  $h + 1$  will be specified simultaneously and if all the values above diagonal of Matrix  $D$  are non-zero, the

algorithm stops.

In [21] theorem 2.2, to calculate  $Sz^*(G)$ , a formula has been presented based on the number of vertices, the number of edges,  $N_u$ ,  $N_v$  and  $Sz(G)$ . To use this formula, we need to calculate  $N_u$ ,  $N_v$  and  $Sz(G)$ . Additionally, in corollary 2.3, for values  $Sz(G)$  and  $Sz^*(G)$ , an upper bound has been identified on the basis of the number of edges and vertices of the graph. In this paper, we will present a simple and applicable algorithm for calculating  $N_u$ ,  $N_v$ ,  $Sz(G)$  and  $Sz^*(G)$  based on distance Matrix  $D$ .

Now, with the help of Matrix  $D$  and the following algorithm, we can simply calculate  $Sz(G)$  and  $Sz^*(G)$ .

In [8], values  $Sz(G)$  and  $Sz^*(G)$  for some specific graphs have been achieved based on values of Wiener index. However, for calculating Wiener index, we need to identify  $d(u, v)$  (the number of edges on any of the shortest paths joining vertex  $u$  to vertex  $v$ ) and in this papers, we have presented the method of calculating  $d(u, v)$ , [8] has not presented a method for calculating it.

Algorithm for calculating  $Sz(G)$  and  $Sz^*(G)$ :

Let

$$Sz = 0, \quad Sz^* = 0, \quad n_i = 0, \quad n_j = 0, \quad n_0 = 0.$$

For all the values above diagonal of Matrix  $D$  with  $D(i, j) = 1$ , (This means that there is an edge  $e$  between vertices  $i$  and  $j$ ):

1. for  $k = 1 : n$   
calculate  $D(i, k) - D(j, k)$ ,
2. If  $D(i, k) - D(j, k) < 0$ ,  
 $n_i = n_i + 1$ ,  
(Vertices whose distance from vertex  $v_j$  is greater than their distance from vertex  $v_i$ )
3. If  $D(i, k) - D(j, k) > 0$ ,  
 $n_j = n_j + 1$ ,
4. If  $D(i, k) - D(j, k) = 0$ ,  
 $n_0 = n_0 + 1$ ,

We do not need Step 3 to calculate  $Sz$ .

$$\begin{cases} Sz^* = Sz^* + \left(n_i + \frac{n_0}{2}\right) \cdot \left(n_j + \frac{n_0}{2}\right); \\ Sz = Sz + (n_i) \cdot (n_j). \end{cases}$$

### 3. NUMERICAL EXAMPLES

In this part, some examples are given to study the efficiency of the mentioned method and also to compare it with other solving methods [8].

Example 1: Table (1) shows the values of  $Sz(G)$  and  $Sz^*(G)$  for some graphs, respectively.

The values obtained for Complete graphs,

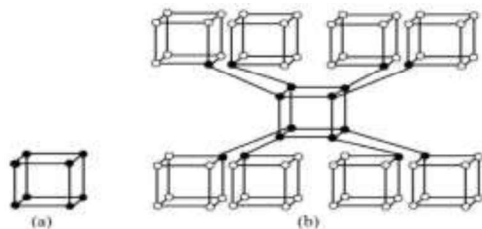
Complete bipartite graphs and Path graphs are the same as the results found in [8].

The values obtained for Fullerene graphs are the same as those of [22]. Therefore, this algorithm is practical for calculating  $Sz(G)$  and  $Sz^*(G)$ .

**Table 1.** The values of  $Sz(G)$  and  $Sz^*(G)$  for special graphs.

G	$Sz(G)$	$Sz^*(G)$
Complete graphs $K_n$	$\frac{n(n-1)}{2}$	$\frac{n^3(n-1)}{8}$
$K_7$	21	257.25
$K_8$	28	448.00
$K_9$	36	729.00
Complete bipartite graphs $K_{n,n}$	$n^4$	$n^4$
$K_{7,7}$	2401	2401
$K_{8,8}$	4096	4096
$K_{9,9}$	6561	6561
Path graphs $P_n$	$\frac{n^3-n}{6}$	$\frac{n^3-n}{6}$
$P_8$	84	84
$P_9$	120	120
$P_{10}$	165	165
Fullerene graphs $C_n$		
$C_{72}$	95412	139188
$C_{84}$	163896	222024
$C_{96}$	235236	331260

Example 2. Let  $G = CCC(n)$  be the crystal structure of cubic carbon. values of  $Sz(G)$  and  $Sz^*(G)$  of  $CCC(n)$  for  $n = 1$  and  $n = 2$  is equal to  $Sz(G) = Sz^*(G) = 192$  and  $Sz(G) = Sz^*(G) = 4.0704 \times 10^4$ . The values obtained for Crystal Structure Cubic Carbon are the same as those of [23].



**Fig. 1.** Crystal Structure Cubic Carbon CCC (1) in part (a) and CCC (2) in part (b).

#### 4. CONCLUSION

Distance Matrix  $D$  has been used to generate a number of topological indices. Many studies have been done to find  $Sz(G)$  and  $Sz^*(G)$  for different graphs. In most cases, the upper and lower bounds are specified for these values. In some graphs, such as Complete graphs and paths, these values are calculated in terms of the number of vertices. In this paper, we calculated  $D$  using adjacency matrix and without multiplying the of matrix. Then, using Matrix  $D$ , the values of  $Sz(G)$  and  $Sz^*(G)$  are easily obtained. The results are completely

acceptable and satisfactory to compared the mentioned references.

#### REFERENCES

- [1] M. Aouchiche and P. Hansena, On a conjecture about the Szeged index, *Eur. J. Comb.* 31 (2010) 1662–1666.
- [2] P.V. Khadikar, N.V. Kale, P.P. Dobrynin, I. Gutman and G. Domotor, The Szeged index and an analogy with the Wiener index, *J. Chem. Inform. Comput. sci.*, 35 (1995) 547–550.
- [3] S. Klavzar, A. Rajapakse and I. Gutman, The Szeged and the Wiener index of graphs, *App. Math. Lett.*, 9 (1996) 45–49.
- [4] T. Pisanski and M. Randic, Use of the Szeged index and the revised Szeged index for measuring network bipartivity, *Dis. App. Math.*, 158 (2010) 1936–1944.
- [5] M. Randic, On generalization of Wiener index for cyclic structures, *Acta Chim. Slovenica*, 49 (2002) 483–496.
- [6] K.C. Das and I. Gutman, Estimating the Szeged index, *App. Math. Lett.*, 22 (2009) 1680–1684.
- [7] X.L. Li and M.M. Liu, Bicyclic graphs with maximal revised Szeged index, *Disc. App. Math.*, 161 (2013) 2527–2531.
- [8] T. Pisanski and J. Zerovnik, Edge contributions of some topological indices and arboreality of molecular graphs, *Ars Math. Contemporanea*, 2 (2009) 49–58.
- [9] S. Simic and I. Gutman, V. Baltic, Some graphs with extremal Szeged index, *Math. Slovaca*, 50 (2000) 1–15.
- [10] R. Xing and B. Zhou, On the revised Szeged index, *Dis. App. Math.*, 159 (2011) 69–78.
- [11] A.T. Balaban, Topological indices based on topological distances in molecular graphs, *Pure and Applied Chemistry*, 55 (1983) 199–206.
- [12] A.T. Balaban, Highly discriminating distance-based topological index, *Chem. Phys. Lett.*, 89 (1989) 399–404.
- [13] H. Hosoya, Topological index: A newly proposed quantity characterizing the topological nature of structural isomers of saturated hydrocarbons, *Bull. Chem. Soc. Japan*, 44 (1971) 2332–2339.
- [14] R. Todeschini and V. Consonni, *Handbook of Molecular Descriptors*, Wiley-VCH, Weinheim, 2000.
- [15] I. Gutman, W. Linert, I. Lukovits and Z. Tomovic, On the multiplicative Wiener index and its possible chemical applications, *Monatshefte fur Chemie*, 131 (2000) 421–427.
- [16] I. Gutman, W. Linert, I. Lukovits and Z. Tomovic, The multiplicative version of the

- Wiener index, *J. Chem. Inf. Comput. Sci.*, 40 (2000) 113–116.
- [17] B. Lucic, I. Lukovits, S. Nikolic and N. Trinajstic, Distance-related indexes in the quantitative structure-property relationship modeling, *J. Chem. Inf. Comput. Sci.*, 41 (2001) 527–535.
- [18] P.G. Seybold, Topological influences on the carcinogenicity of aromatic hydrocarbons. I. Bay region geometry, *Int. J. Quantum Chem.*, 10 (1983) 95–101.
- [19] W.R. Muller, K. Szymanski, J.V. Knop and N. Trinajstic, An Algorithm for construction of the molecular distance matrix, *J. Comput. Chemistry*, 8 (1987) 170–173.
- [20] <https://stackoverflow.com/questions/32164012/how-to-get-distance-matrix-from-adjacency-matrix-matlab>, May, 2021.
- [21] M. Faghani, A.R. Ashraf, Revised and edge revised Szeged indices of graphs, *Ars Math. Contemporanea*, 7 (2013) 153–160.
- [22] A. Mottaghi and Z. Mehranian, PI, Szeged and Revised Szeged Indices of IPR Fullerenes, *Iranian J. Math. Chem.*, 2 (2) (2011) 87-99.
- [23] H. Yang, M. Naeem, A. Q. Baig, H. Shaker and M. K. Siddiqui, Vertex Szeged index of crystal cubic carbon structure, *J. Dis. Math. Sci. Crypto.*, 22 (2019) 1177–1187.

## COPYRIGHTS



© 2022 by the authors. Licensee PNU, Tehran, Iran. This article is an open access article distributed under the terms and conditions of the Creative Commons Attribution 4.0 International (CC BY4.0) (<http://creativecommons.org/licenses/by/4.0>)

## محاسبه شاخص Szeged و Szeged اصلاح شده با استفاده از ماتریس مجاورت

هاجر علیمراد

بخش ریاضی، دانشگاه جهرم، ایران

تاریخ دریافت: ۱۸ تیر ۱۴۰۱ تاریخ پذیرش: ۲۴ مرداد ۱۴۰۱

### چکیده

شاخص‌های توپولوژیک مقادیر ثابت مولکولی هستند که در شیمی نظری برای شناسایی طراحی ترکیبات شیمیایی مولکول‌ها با ویژگی‌های فیزیکی شیمیایی داده شده یا فعالیت‌های دارویی و بیولوژیکی معین استفاده می‌شوند. شاخص Szeged ( $Sz(G)$ ) و Szeged اصلاح شده ( $Sz^*(G)$ ) از خصوصیات فاصله را برای نمودارها مشخص می‌کند. در شیمی محاسباتی و نظریه گراف،  $Sz(G)$  و  $Sz^*(G)$  برای تعیین ویژگی‌های ساختارهای مولکولی شیمیایی بیشتر کاربرد دارند و بنابراین به طور گسترده در کاربردهای شیمیایی مورد استفاده قرار می‌گیرند. در این مقاله، یک الگوریتم ساده برای ایجاد ماتریس فاصله ارائه شده است. این الگوریتم برای محاسبه  $Sz(G)$  و  $Sz^*(G)$  استفاده خواهد شد.

### واژه‌های کلیدی

ماتریس مجاورت، گراف همبند، ماتریس فاصله، شاخص Szeged اصلاح شده، شاخص Szeged.